

Note**On a Theorem of Ryser**

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In this paper we generalize a theorem of Ryser [1] about finite geometries having an equal number n of points and lines, when something is known about the number δ of lines that miss certain lines. The result is that either (a) all lines have k points, all points are on k lines, and each line misses δ' other lines or (b) all lines have at least k points, all points are on at most $k + 1$ lines and n is bounded by a quadratic function of δ . As the bound which is given by the proof is known to be not best possible, we conjecture that $n \leq \delta^2 + 5\delta + 4$ in case (b).

If X is a set of elements (points) and \mathcal{C} is a collection of subsets (blocks or lines) of X , we define:

$$\text{for all } x \in X, \quad r_x = |\{A \in \mathcal{C} : x \in A\}|;$$

$$\text{for all } A \in \mathcal{C}, \quad \delta_A = |\{B \in \mathcal{C} : A \cap B = \emptyset\}|;$$

and

$$k = \min_{A \in \mathcal{C}} |A|.$$

Throughout we will also make use of the duality between blocks and points to write $A \in x$ to mean A is a block through, containing, or on x just as we would write $x \in A$ to mean x is a point on A .

THEOREM 1. *Let X be a set of n points, \mathcal{C} a collection of n blocks on X , and δ a positive integer such that:*

- (i) for all $A \in \mathcal{C}$, $|A| \geq 3$;
- (ii) for all $A, B \in \mathcal{C}$, $A \neq B$, $|A \cap B| \leq 1$; and
- (iii) $|A| = k$ implies $\delta_A \leq \delta \leq k - 2$.

Then either

(a) for all $x \in X$, $r_x = k$, for all $A \in \mathcal{C}$, $|A| = k$, and for all $A \in \mathcal{C}$, $\delta_A = \delta' \leq \delta$, or

(b) for all $x \in X$, $r_x \leq k + 1$, for all $A \in \mathcal{C}$, $|A| \geq k$, and $n = k^2 + \epsilon$, where $k + \epsilon \leq 2\delta + 2$ (which forces $n \leq (2\delta + 2)^2$).

Proof. We first state the three basic counting arguments to be used in the proof. Suppose $x \in X$ and $A, B \in \mathcal{C}$ with $\{x\} = A \cap B$.

Counting points on the blocks through x ,

$$n = |X| \geq 1 + \sum_{D \ni x} (|D| - 1). \quad (1)$$

Counting blocks missing A and blocks through points on A ,

$$n = |\mathcal{C}| = \delta_A + \left[1 + \sum_{y \in A} (r_y - 1) \right]. \quad (2)$$

And counting blocks missing A or B and blocks intersecting both,

$$n = |\mathcal{C}| \leq (\delta_A + \delta_B) + [r_x + (|A| - 1)(|B| - 1)]. \quad (3)$$

Now fix $A \in \mathcal{C}$ with $|A| = k$, and choose $x \in A$ such that $r_x = \max_{y \in A} \{r_y\}$. If x were on no other blocks of size k , then from (1) and (2)

$$1 + 1 \cdot (k - 1) + (r_x - 1)k \leq n \leq 1 + \delta_A + k(r_x - 1).$$

So

$$k \leq 1 + \delta_A,$$

which contradicts hypothesis (iii). Thus there exists a second block B through x with $|B| = k$.

Now applying (1) and (3), we get

$$1 + r_x(k - 1) \leq n \leq \delta_A + \delta_B + r_x + (k - 1)^2 \quad (4)$$

so

$$r_x \leq k + [(\delta_A + \delta_B)/(k - 2)] \leq k + 2$$

since $\delta_A, \delta_B \leq k - 2$ by (iii).

Suppose for the moment that $r_y \leq k$ for all y on at least one block of size k . If z is a point which is not on any block of size k , then from (1) and (2),

$$1 + r_z \cdot k \leq n \leq 1 + \delta_A + k(k - 1),$$

$$r_z \leq k - 1 + (\delta_A/k) < k.$$

Counting point-block incidences two ways,

$$n \cdot k \leq \sum_{A \in \mathcal{O}} |A| = \sum_{x \in X} r_x \leq n \cdot k.$$

So equality holds, giving $|A| = k$ for all $A \in \mathcal{O}$. And from (2), $n = 1 + \delta_A + k(k-1)$ for all $A \in \mathcal{O}$, so $\delta_A = \delta' \leq \delta$ for all $A \in \mathcal{O}$ and conclusion (a) holds.

Returning to (4), suppose $r_x = k+2$. This forces $\delta_A = \delta_B = \delta = k-2$ and

$$1 + (k+2)(k-1) = n = (k-2) + (k-2) + (k+2) + (k-1)^2. \quad (4')$$

This means all blocks through x have size k ; all points must be on some block through x , for all $y \neq x$; $r_y \geq 1 + (k-1) = k$, since y is on a block through x of size k ; and $|D| \geq k+1$ for all $D \in \mathcal{O}$, D not on x , since no block missing x can miss two blocks on x . Thus if $y \neq x$, $r_y \neq k+2$, and if $r_y = k+1$, then from (1) and (4')

$$1 + 1 \cdot (k-1) + k \cdot k \leq n = 1 + (k+2)(k-1),$$

which is a contradiction. So $r_y = k$ for all $y \neq x$. Counting point-block incidences,

$$\begin{aligned} (n-1)k + 1(k-2) &= \sum_{x \in X} r_x = \sum_{A \in \mathcal{O}} |A| \\ &\geq (k+2)k + (n-k-2)(k+1). \end{aligned}$$

So

$$\begin{aligned} 2k+4 &\geq n = 1 + (k+2)(k-1), \\ k^2 - k - 5 &\leq 0. \end{aligned}$$

But from (i), $k \geq 3$. So $r_x \neq k+2$.

We now have $r_y \leq k+1$ for all y on at least one block of size k . If z is not on any block of size k , then from (1) and (2)

$$\begin{aligned} 1 + r_z k &\leq n \leq 1 + \delta + k \cdot k, \\ r_z &\leq k + (\delta/k) < k+1. \end{aligned}$$

So $r_y \leq k+1$ for all $y \in X$. From (1)

$$n \geq 1 + (k+1)(k-1) = k^2, \quad \text{so let } n = k^2 + \epsilon.$$

Then from (3),

$$k^2 + \epsilon = n \leq 2\delta + (k+1) + (k-1)^2$$

giving $k + \epsilon \leq 2\delta + 2$, and (b) holds, proving the theorem.

The bound obtained in the proof is not best possible. It is conjectured that the best bound is $n \leq \delta^2 + 5\delta + 4$, which conforms to the known results of Ryser [1] and Keenan [unpublished] for $\delta = 1$ and $\delta = 2$, respectively. One would hope that the simple geometric construction for the maximal sizes would generalize past $\delta = 2$, and to prove the bound one need only show $\delta_A = k - 2$ for some A with $|A| = k$. It would be nice if these special configurations were self-dual.

The hypothesis that $\delta \leq k - 2$ is not completely necessary, so we prove the following.

THEOREM 2. *Replacing (iii) in Theorem 1 by*

$$(iii') \quad \delta_A \leq \delta \text{ for all } A \in \mathcal{C} \text{ with } |A| = k \text{ or } k + 1.$$

Then either Theorem 1 applies or

$$(c) \quad k \leq \delta + 1 \text{ and } n \leq \delta^2 + 5\delta + 2.$$

Proof. If $k \geq \delta + 2$, then Theorem 1 applies, so we may assume $k \leq \delta + 1$. Let $r = \max_{x \in X} \{r_x\}$.

Suppose two blocks of size k intersect. Then as before from (1) and (3)

$$\begin{aligned} 1 + r(k - 1) &\leq n \leq 2\delta + r + (k - 1)^2 \\ r &\leq k + [2\delta/(k - 2)] \leq k + 2\delta. \end{aligned}$$

And since $k \leq \delta + 1$,

$$\begin{aligned} n &\leq 2\delta + (k + 2\delta) + (k - 1)^2, \\ n &\leq 2\delta + (3\delta + 1) + \delta^2 = \delta^2 + 5\delta + 1. \end{aligned}$$

If A and B are two blocks of size k which do not intersect, then by (2)

$$\begin{aligned} n &\leq \delta_B + (k \cdot k + \delta_A), \\ n &\leq 2\delta + (\delta + 1)^2 = \delta^2 + 4\delta + 1. \end{aligned}$$

So we may assume that there is only one block A of size k . Let B be the next smallest block, and let $|B| = k + \alpha$. From (1) and (2),

$$\begin{aligned} 1 + (r - 1)(k + \alpha - 1) + 1 \cdot (k - 1) &\leq n \leq 1 + \delta + k(r - 1), \\ (r - 1)(\alpha - 1) &\leq \delta + 1 - k. \end{aligned}$$

If $\alpha > 1$

$$r \leq \delta + 1 - k \leq \delta - 2$$

and

$$n \leq 1 + \delta + (\delta + 1)(\delta - 3) = \delta^2 - \delta - 2.$$

So $|B| = k + 1$.

If $A \cap B \neq \emptyset$, then from (1) and (3)

$$1 + (r - 1)k + 1 \cdot (k - 1) \leq n \leq 2\delta + r + k(k - 1),$$

$$r \leq k + [2\delta/(k - 1)] \leq k + \delta.$$

So

$$n \leq 2\delta + (k + \delta) + k(k - 1),$$

$$n \leq 2\delta + (2\delta + 1) + (\delta + 1)\delta = \delta^2 + 5\delta + 1.$$

If $A \cap B = \emptyset$, then from (2)

$$n \leq \delta_B + (k(k + 1) + \delta_A),$$

$$n \leq 2\delta + (\delta + 1)(\delta + 2) = \delta^2 + 5\delta + 2,$$

which finishes the proof.

REFERENCE

1. H. J. RYSER, Subsets of a finite set that intersect each other in at most one element, *J. Combinatorial Theory A* **17** (1974), 59-74.